Data Representation

Chapter One

Probably the biggest stumbling block most beginners encounter when attempting to learn assembly language is the common use of the binary and hexadecimal numbering systems. Many programmers think that hexadecimal (or hex\(^1\)) numbers represent absolute proof that God never intended anyone to work in assembly language. While it is true that hexadecimal numbers are a little different from what you may be used to, their advantages outweigh their disadvantages by a large margin. Nevertheless, understanding these numbering systems is important because their use simplifies other complex topics including boolean algebra and logic design, signed numeric representation, character codes, and packed data.

1.0 Chapter Overview

This chapter discusses several important concepts including the binary and hexadecimal numbering systems, binary data organization (bits, nibbles, bytes, words, and double words), signed and unsigned numbering systems, arithmetic, logical, shift, and rotate operations on binary values, bit fields and packed data, and the ASCII character set. This is basic material and the remainder of this text depends upon your understanding of these concepts. If you are already familiar with these terms from other courses or study, you should at least skim this material before proceeding to the next chapter. If you are unfamiliar with this material, or only vaguely familiar with it, you should study it carefully before proceeding. \textit{All of the material in this chapter is important!} Do not skip over any material.

1.1 Numbering Systems

Most modern computer systems do not represent numeric values using the decimal system. Instead, they typically use a binary or two’s complement numbering system. To understand the limitations of computer arithmetic, you must understand how computers represent numbers.

1.1.1 A Review of the Decimal System

You’ve been using the decimal (base 10) numbering system for so long that you probably take it for granted. When you see a number like “123”, you don’t think about the value 123; rather, you generate a mental image of how many items this value represents. In reality, however, the number 123 represents:

\[ 1 \times 10^2 + 2 \times 10^1 + 3 \times 10^0 \]

or

\[ 100 + 20 + 3 \]

Each digit appearing to the left of the decimal point represents a value between zero and nine times an increasing power of ten. Digits appearing to the right of the decimal point represent a value between zero and nine times an increasing negative power of ten. For example, the value 123.456 means:

\[ 1 \times 10^2 + 2 \times 10^1 + 3 \times 10^0 + 4 \times 10^{-1} + 5 \times 10^{-2} + 6 \times 10^{-3} \]

or

\[ 100 + 20 + 3 + 0.4 + 0.05 + 0.006 \]

1. Hexadecimal is often abbreviated as \textit{hex} even though, technically speaking, \textit{hex} means base six, not base sixteen.
1.1.2 The Binary Numbering System

Most modern computer systems (including the IBM PC) operate using binary logic. The computer represents values using two voltage levels (usually 0v and +5v). With two such levels we can represent exactly two different values. These could be any two different values, but by convention we use the values zero and one. These two values, coincidentally, correspond to the two digits used by the binary numbering system. Since there is a correspondence between the logic levels used by the 80x86 and the two digits used in the binary numbering system, it should come as no surprise that the IBM PC employs the binary numbering system.

The binary numbering system works just like the decimal numbering system, with two exceptions: binary only allows the digits 0 and 1 (rather than 0-9), and binary uses powers of two rather than powers of ten. Therefore, it is very easy to convert a binary number to decimal. For each “1” in the binary string, add in $2^n$ where “n” is the zero-based position of the binary digit. For example, the binary value $11001010_2$ represents:

$$1 \cdot 2^7 + 1 \cdot 2^6 + 0 \cdot 2^5 + 0 \cdot 2^4 + 1 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 0 \cdot 2^0 = 128 + 64 + 8 + 2 = 202_{10}$$

To convert decimal to binary is slightly more difficult. You must find those powers of two which, when added together, produce the decimal result. The easiest method is to work from the a large power of two down to $2^0$. Consider the decimal value 1359:

- $2^{10} = 1024, 2^{11} = 2048$. So 1024 is the largest power of two less than 1359. Subtract 1024 from 1359 and begin the binary value on the left with a “1” digit. Binary = “1”, Decimal result is 1359 - 1024 = 335.
- The next lower power of two ($2^9 = 512$) is greater than the result from above, so add a “0” to the end of the binary string. Binary = “10”, Decimal result is still 335.
- The next lower power of two is 256 ($2^8$). Subtract this from 335 and add a “1” digit to the end of the binary number. Binary = “101”, Decimal result is 79.
- 128 ($2^7$) is greater than 79, so tack a “0” to the end of the binary string. Binary = “1010”, Decimal result remains 79.
- The next lower power of two is 256 ($2^8$). Subtract this from 335 and add a “1” digit to the end of the binary number. Binary = “10101”, Decimal result is 15.
- 15 is less than the next power of two ($2^5 = 32$) so simply add a “0” to the end of the binary string. Binary = “101010”, Decimal result is still 15.
- 16 ($2^4$) is greater than the remainder so far, so append a “0” to the end of the binary string. Binary = “1010100”, Decimal result is 15.
- $2^3$ (eight) is less than 15, so stick another “1” digit on the end of the binary string. Binary = “10101001”, Decimal result is 7.
- $2^2$ is less than seven, so subtract four from seven and append another one to the binary string. Binary = “101010011”, decimal result is 3.
- $2^1$ is less than three, so append a one to the end of the binary string and subtract two from the decimal value. Binary = “1010100111”, Decimal result is now 1.
- Finally, the decimal result is one, which is $2^0$, so add a final “1” to the end of the binary string. The final binary result is “10101001111”
Binary numbers, although they have little importance in high level languages, appear everywhere in assembly language programs.

### 1.1.3 Binary Formats

In the purest sense, every binary number contains an infinite number of digits (or bits which is short for binary digits). For example, we can represent the number five by:

\[
\begin{array}{ccccccccccccccc}
101 & 00000101 & 000000000101 & \ldots & \phantom{000000000000101}
\end{array}
\]

Any number of leading zero bits may precede the binary number without changing its value.

We will adopt the convention ignoring any leading zeros. For example, 101₂ represents the number five. Since the 80x86 works with groups of eight bits, we’ll find it much easier to zero extend all binary numbers to some multiple of four or eight bits. Therefore, following this convention, we’d represent the number five as 0101₂ or 00000101₂.

In the United States, most people separate every three digits with a comma to make larger numbers easier to read. For example, 1,023,435,208 is much easier to read and comprehend than 1023435208. We’ll adopt a similar convention in this text for binary numbers. We will separate each group of four binary bits with a space. For example, the binary value 1010111110110010 will be written 1010 1111 1011 0010.

We often pack several values together into the same binary number. One form of the 80x86 MOV instruction (see appendix D) uses the binary encoding 1011 0rrr dddd dddd to pack three items into 16 bits: a five-bit operation code (10110), a three-bit register field (rrr), and an eight-bit immediate value (dddd dddd). For convenience, we’ll assign a numeric value to each bit position. We’ll number each bit as follows:

1) The rightmost bit in a binary number is bit position zero.
2) Each bit to the left is given the next successive bit number.

An eight-bit binary value uses bits zero through seven:

\[
x_7 \ x_6 \ x_5 \ x_4 \ x_3 \ x_2 \ x_1 \ x_0
\]

A 16-bit binary value uses bit positions zero through fifteen:

\[
x_{15} \ x_{14} \ x_{13} \ x_{12} \ x_{11} \ x_{10} \ x_9 \ x_8 \ x_7 \ x_6 \ x_5 \ x_4 \ x_3 \ x_2 \ x_1 \ x_0
\]

Bit zero is usually referred to as the low order (L.O.) bit. The left-most bit is typically called the high order (H.O.) bit. We’ll refer to the intermediate bits by their respective bit numbers.

### 1.2 Data Organization

In pure mathematics a value may take an arbitrary number of bits. Computers, on the other hand, generally work with some specific number of bits. Common collections are single bits, groups of four bits (called nibbles), groups of eight bits (called bytes), groups of 16 bits (called words), and more. The sizes are not arbitrary. There is a good reason for these particular values. This section will describe the bit groups commonly used on the Intel 80x86 chips.
1.2.1 Bits

The smallest “unit” of data on a binary computer is a single bit. Since a single bit is capable of representing only two different values (typically zero or one) you may get the impression that there are a very small number of items you can represent with a single bit. Not true! There are an infinite number of items you can represent with a single bit.

With a single bit, you can represent any two distinct items. Examples include zero or one, true or false, on or off, male or female, and right or wrong. However, you are not limited to representing binary data types (that is, those objects which have only two distinct values). You could use a single bit to represent the numbers 723 and 1,245. Or perhaps 6,254 and 5. You could also use a single bit to represent the colors red and blue. You could even represent two unrelated objects with a single bit. For example, you could represent the color red and the number 3,256 with a single bit. You can represent any two different values with a single bit. However, you can represent only two different values with a single bit.

To confuse things even more, different bits can represent different things. For example, one bit might be used to represent the values zero and one, while an adjacent bit might be used to represent the values true and false. How can you tell by looking at the bits? The answer, of course, is that you can’t. But this illustrates the whole idea behind computer data structures: data is what you define it to be. If you use a bit to represent a boolean (true/false) value then that bit (by your definition) represents true or false. For the bit to have any true meaning, you must be consistent. That is, if you’re using a bit to represent true or false at one point in your program, you shouldn’t use the true/false value stored in that bit to represent red or blue later.

Since most items you’ll be trying to model require more than two different values, single bit values aren’t the most popular data type you’ll use. However, since everything else consists of groups of bits, bits will play an important role in your programs. Of course, there are several data types that require two distinct values, so it would seem that bits are important by themselves. However, you will soon see that individual bits are difficult to manipulate, so we’ll often use other data types to represent boolean values.

1.2.2 Nibbles

A nibble is a collection of four bits. It wouldn’t be a particularly interesting data structure except for two items: BCD (binary coded decimal) numbers and hexadecimal numbers. It takes four bits to represent a single BCD or hexadecimal digit. With a nibble, we can represent up to 16 distinct values. In the case of hexadecimal numbers, the values 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, and F are represented with four bits (see “The Hexadecimal Numbering System” on page 17). BCD uses ten different digits (0, 1, 2, 3, 4, 5, 6, 7, 8, 9) and requires four bits. In fact, any sixteen distinct values can be represented with a nibble, but hexadecimal and BCD digits are the primary items we can represent with a single nibble.

1.2.3 Bytes

Without question, the most important data structure used by the 80x86 microprocessor is the byte. A byte consists of eight bits and is the smallest addressable datum (data item) on the 80x86 microprocessor. Main memory and I/O addresses on the 80x86 are all byte addresses. This means that the smallest item that can be individually accessed by an 80x86 program is an eight-bit value. To access anything smaller requires that you read the byte containing the data and mask out the unwanted bits. The bits in a byte are normally numbered from zero to seven using the convention in Figure 1.1.

Bit 0 is the low order bit or least significant bit, bit 7 is the high order bit or most significant bit of the byte. We’ll refer to all other bits by their number.
Note that a byte also contains exactly two nibbles (see Figure 1.2).

![Figure 1.1: Bit Numbering in a Byte](image)

Bits 0..3 comprise the low order nibble, bits 4..7 form the high order nibble. Since a byte contains exactly two nibbles, byte values require two hexadecimal digits.

Since a byte contains eight bits, it can represent $2^8$, or 256, different values. Generally, we’ll use a byte to represent numeric values in the range 0..255, signed numbers in the range -128..+127 (see “Signed and Unsigned Numbers” on page 23), ASCII/IBM character codes, and other special data types requiring no more than 256 different values. Many data types have fewer than 256 items so eight bits is usually sufficient.

Since the 80x86 is a byte addressable machine (see “Memory Layout and Access” on page 145), it turns out to be more efficient to manipulate a whole byte than an individual bit or nibble. For this reason, most programmers use a whole byte to represent data types that require no more than 256 items, even if fewer than eight bits would suffice. For example, we’ll often represent the boolean values true and false by $00000001_2$ and $00000000_2$ (respectively).

Probably the most important use for a byte is holding a character code. Characters typed at the keyboard, displayed on the screen, and printed on the printer all have numeric values. To allow it to communicate with the rest of the world, the IBM PC uses a variant of the ASCII character set (see “The ASCII Character Set” on page 28). There are 128 defined codes in the ASCII character set. IBM uses the remaining 128 possible values for extended character codes including European characters, graphic symbols, Greek letters, and math symbols. See Appendix A for the character/code assignments.

### 1.2.4 Words

A word is a group of 16 bits. We’ll number the bits in a word starting from zero on up to fifteen. The bit numbering appears in Figure 1.3.

![Figure 1.3: Bit Numbers in a Word](image)

Like the byte, bit 0 is the low order bit and bit 15 is the high order bit. When referencing the other bits in a word use their bit position number.
Notice that a word contains exactly two bytes. Bits 0 through 7 form the low order byte, bits 8 through 15 form the high order byte (see Figure 1.4).

![Figure 1.4: The Two Bytes in a Word](image)

Naturally, a word may be further broken down into four nibbles as shown in Figure 1.5.

![Figure 1.5: Nibbles in a Word](image)

Nibble zero is the low order nibble in the word and nibble three is the high order nibble of the word. The other two nibbles are “nibble one” or “nibble two”.

With 16 bits, you can represent \(2^{16}\) (65,536) different values. These could be the values in the range 0..65,535 (or, as is usually the case, -32,768..+32,767) or any other data type with no more than 65,536 values. The three major uses for words are integer values, offsets, and segment values (see “Memory Layout and Access” on page 145 for a description of segments and offsets).

Words can represent integer values in the range 0..65,535 or -32,768..32,767. Unsigned numeric values are represented by the binary value corresponding to the bits in the word. Signed numeric values use the two’s complement form for numeric values (see “Signed and Unsigned Numbers” on page 23). Segment values, which are always 16 bits long, constitute the paragraph address of a code, data, extra, or stack segment in memory.

### 1.2.5 Double Words

A double word is exactly what its name implies, a pair of words. Therefore, a double word quantity is 32 bits long as shown in Figure 1.6.

![Figure 1.6: Bit Numbers in a Double Word](image)

Naturally, this double word can be divided into a high order word and a low order word, or four different bytes, or eight different nibbles (see Figure 1.7).

Double words can represent all kinds of different things. First and foremost on the list is a segmented address. Another common item represented with a double word is a 32-bit...
integer value (which allows unsigned numbers in the range 0..4,294,967,295 or signed numbers in the range -2,147,483,648..2,147,483,647). 32-bit floating point values also fit into a double word. Most of the time, we’ll use double words to hold segmented addresses.

1.3 The Hexadecimal Numbering System

A big problem with the binary system is verbosity. To represent the value 202\textsubscript{10} requires eight binary digits. The decimal version requires only three decimal digits and, thus, represents numbers much more compactly than does the binary numbering system. This fact was not lost on the engineers who designed binary computer systems. When dealing with large values, binary numbers quickly become too unwieldy. Unfortunately, the computer thinks in binary, so most of the time it is convenient to use the binary numbering system. Although we can convert between decimal and binary, the conversion is not a trivial task. The hexadecimal (base 16) numbering system solves these problems. Hexadecimal numbers offer the two features we’re looking for: they’re very compact, and it’s simple to convert them to binary and vice versa. Because of this, most binary computer systems today use the hexadecimal numbering system\textsuperscript{2}. Since the radix (base) of a hexadecimal number is 16, each hexadecimal digit to the left of the hexadecimal point represents some value times a successive power of 16. For example, the number 1234\textsubscript{16} is equal to:

\[
1 \times 16^3 + 2 \times 16^2 + 3 \times 16^1 + 4 \times 16^0
\]

or

\[
4096 + 512 + 48 + 4 = 4660\textsubscript{10}.
\]

Each hexadecimal digit can represent one of sixteen values between 0 and 15\textsubscript{10}. Since there are only ten decimal digits, we need to invent six additional digits to represent the values in the range 10\textsubscript{10} through 15\textsubscript{10}. Rather than create new symbols for these digits, we’ll use the letters A through F. The following are all examples of valid hexadecimal numbers:

\[
\text{Figure 1.7: Nibbles, Bytes, and Words in a Double Word}
\]

\[
\begin{array}{cccccc}
31 & 23 & 15 & 7 & 0 \\
\text{H.O. Word} & \text{L.O. Word} \\
\hline
31 & 23 & 15 & 7 & 0 \\
\text{H.O. Byte} & \text{Byte # 2} & \text{Byte # 1} & \text{L.O. Byte} \\
\hline
31 & 23 & 15 & 7 & 0 \\
\text{Nibble #7} & \#6 & \#5 & \#4 & \#3 & \#2 & \#1 & \#0 \\
\text{H. O.} & \text{L. O.}
\end{array}
\]

\[\text{Figure 1.7: Nibbles, Bytes, and Words in a Double Word}\]

\[2. \text{Digital Equipment is the only major holdout. They still use octal numbers in most of their systems. A legacy of the days when they produced 12-bit machines.}\]
1234_{16} \ DEAD_{16} \ BEEF_{16} \ 0AFB_{16} \ FEED_{16} \ DEAF_{16}

Since we’ll often need to enter hexadecimal numbers into the computer system, we’ll need a different mechanism for representing hexadecimal numbers. After all, on most computer systems you cannot enter a subscript to denote the radix of the associated value. We’ll adopt the following conventions:

- All numeric values (regardless of their radix) begin with a decimal digit.
- All hexadecimal values end with the letter “h”, e.g., 123A4h³.
- All binary values end with the letter “b”.
- Decimal numbers may have a “t” or “d” suffix.

Examples of valid hexadecimal numbers:

1234h \ 0DEADh \ 0BEEFh \ 0AFBh \ 0FEEDh \ 0DEAh

As you can see, hexadecimal numbers are compact and easy to read. In addition, you can easily convert between hexadecimal and binary. Consider the following table:

<table>
<thead>
<tr>
<th>Binary</th>
<th>Hexadecimal</th>
</tr>
</thead>
<tbody>
<tr>
<td>0000</td>
<td>0</td>
</tr>
<tr>
<td>0001</td>
<td>1</td>
</tr>
<tr>
<td>0010</td>
<td>2</td>
</tr>
<tr>
<td>0011</td>
<td>3</td>
</tr>
<tr>
<td>0100</td>
<td>4</td>
</tr>
<tr>
<td>0101</td>
<td>5</td>
</tr>
<tr>
<td>0110</td>
<td>6</td>
</tr>
<tr>
<td>0111</td>
<td>7</td>
</tr>
<tr>
<td>1000</td>
<td>8</td>
</tr>
<tr>
<td>1001</td>
<td>9</td>
</tr>
<tr>
<td>1010</td>
<td>A</td>
</tr>
<tr>
<td>1011</td>
<td>B</td>
</tr>
<tr>
<td>1100</td>
<td>C</td>
</tr>
<tr>
<td>1101</td>
<td>D</td>
</tr>
<tr>
<td>1110</td>
<td>E</td>
</tr>
<tr>
<td>1111</td>
<td>F</td>
</tr>
</tbody>
</table>

This table provides all the information you’ll ever need to convert any hexadecimal number into a binary number or vice versa.

To convert a hexadecimal number into a binary number, simply substitute the corresponding four bits for each hexadecimal digit in the number. For example, to convert

---

3. Actually, following hexadecimal values with an “h” is an Intel convention, not a general convention. The 68000 and 65c816 assemblers used in the Macintosh and Apple II denote hexadecimal numbers by prefacing the hex value with a “$” symbol.
0ABCDh into a binary value, simply convert each hexadecimal digit according to the table above:

<table>
<thead>
<tr>
<th>0</th>
<th>A</th>
<th>B</th>
<th>C</th>
<th>D</th>
<th>Hexadecimal</th>
</tr>
</thead>
<tbody>
<tr>
<td>0000</td>
<td>1010</td>
<td>1011</td>
<td>1100</td>
<td>1101</td>
<td>Binary</td>
</tr>
</tbody>
</table>

To convert a binary number into hexadecimal format is almost as easy. The first step is to pad the binary number with zeros to make sure that there is a multiple of four bits in the number. For example, given the binary number 1011001010, the first step would be to add two bits to the left of the number so that it contains 12 bits. The converted binary value is 001011001010. The next step is to separate the binary value into groups of four bits, e.g., 0010 1100 1010. Finally, look up these binary values in the table above and substitute the appropriate hexadecimal digits, e.g., 2CA. Contrast this with the difficulty of conversion between decimal and binary or decimal and hexadecimal!

Since converting between hexadecimal and binary is an operation you will need to perform over and over again, you should take a few minutes and memorize the table above. Even if you have a calculator that will do the conversion for you, you’ll find manual conversion to be a lot faster and more convenient when converting between binary and hex.

### 1.4 Arithmetic Operations on Binary and Hexadecimal Numbers

There are several operations we can perform on binary and hexadecimal numbers. For example, we can add, subtract, multiply, divide, and perform other arithmetic operations. Although you needn’t become an expert at it, you should be able to, in a pinch, perform these operations manually using a piece of paper and a pencil. Having just said that you should be able to perform these operations manually, the correct way to perform such arithmetic operations is to have a calculator which does them for you. There are several such calculators on the market; the following table lists some of the manufacturers who produce such devices:

Manufacturers of Hexadecimal Calculators:

- Casio
- Hewlett-Packard
- Sharp
- Texas Instruments

This list is, by no means, exhaustive. Other calculator manufacturers probably produce these devices as well. The Hewlett-Packard devices are arguably the best of the bunch. However, they are more expensive than the others. Sharp and Casio produce units which sell for well under $50. If you plan on doing any assembly language programming at all, owning one of these calculators is essential.

Another alternative to purchasing a hexadecimal calculator is to obtain a TSR (Terminate and Stay Resident) program such as SideKick™ which contains a built-in calculator. However, unless you already have one of these programs, or you need some of the other features they offer, such programs are not a particularly good value since they cost more than an actual calculator and are not as convenient to use.

To understand why you should spend the money on a calculator, consider the following arithmetic problem:

\[
\begin{align*}
9h + 1h &= 10h \\
9h &+ 1h \\
\hline
10h
\end{align*}
\]

You’re probably tempted to write in the answer “10h” as the solution to this problem. But that is not correct! The correct answer is ten, which is “0Ah”, not sixteen which is “10h”. A similar problem exists with the arithmetic problem:
You’re probably tempted to answer “9h” even though the true answer is “0Fh”. Remember, this problem is asking “what is the difference between sixteen and one?” The answer, of course, is fifteen which is “0Fh”.

Even if the two problems above don’t bother you, in a stressful situation your brain will switch back into decimal mode while you’re thinking about something else and you’ll produce the incorrect result. Moral of the story – if you must do an arithmetic computation using hexadecimal numbers by hand, take your time and be careful about it. Either that, or convert the numbers to decimal, perform the operation in decimal, and convert them back to hexadecimal.

You should never perform binary arithmetic computations. Since binary numbers usually contain long strings of bits, there is too much of an opportunity for you to make a mistake. Always convert binary numbers to hex, perform the operation in hex (preferably with a hex calculator) and convert the result back to binary, if necessary.

### 1.5 Logical Operations on Bits

There are four main logical operations we’ll need to perform on hexadecimal and binary numbers: AND, OR, XOR (exclusive-or), and NOT. Unlike the arithmetic operations, a hexadecimal calculator isn’t necessary to perform these operations. It is often easier to do them by hand than to use an electronic device to compute them. The logical AND operation is a dyadic\(^4\) operation (meaning it accepts exactly two operands). These operands are single binary (base 2) bits. The AND operation is:

- 0 and 0 = 0
- 0 and 1 = 0
- 1 and 0 = 0
- 1 and 1 = 1

A compact way to represent the logical AND operation is with a truth table. A truth table takes the following form:

<table>
<thead>
<tr>
<th>AND</th>
<th>0</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
</tbody>
</table>

This is just like the multiplication tables you encountered in elementary school. The column on the left and the row at the top represent input values to the AND operation. The value located at the intersection of the row and column (for a particular pair of input values) is the result of logically ANDing those two values together. In English, the logical AND operation is, “If the first operand is one and the second operand is one, the result is one; otherwise the result is zero.”

One important fact to note about the logical AND operation is that you can use it to force a zero result. If one of the operands is zero, the result is always zero regardless of the other operand. In the truth table above, for example, the row labelled with a zero input

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4. Many texts call this a binary operation. The term dyadic means the same thing and avoids the confusion with the binary numbering system.
contains only zeros and the column labelled with a zero only contains zero results. Con-\nversely, if one operand contains a one, the result is exactly the value of the second oper-
and. These features of the AND operation are very important, particularly when working
with bit strings and we want to force individual bits in the string to zero. We will investi-
gate these uses of the logical AND operation in the next section.

The logical OR operation is also a dyadic operation. Its definition is:

\[
\begin{align*}
0 \text{ or } 0 &= 0 \\
0 \text{ or } 1 &= 1 \\
1 \text{ or } 0 &= 1 \\
1 \text{ or } 1 &= 1
\end{align*}
\]

The truth table for the OR operation takes the following form:

**Table 3: OR Truth Table**

<table>
<thead>
<tr>
<th>OR</th>
<th>0</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
</tbody>
</table>

Colloquially, the logical OR operation is, “If the first operand or the second operand
(or both) is one, the result is one; otherwise the result is zero.” This is also known as the
*inclusive-OR* operation.

If one of the operands to the logical-OR operation is a one, the result is always one
regardless of the second operand’s value. If one operand is zero, the result is always the
value of the second operand. Like the logical AND operation, this is an important
side-effect of the logical-OR operation that will prove quite useful when working with bit
strings (see the next section).

Note that there is a difference between this form of the inclusive logical OR operation
and the standard English meaning. Consider the phrase “I am going to the store or I am
going to the park.” Such a statement implies that the speaker is going to the store or to the
park but not to both places. Therefore, the English version of logical OR is slightly differ-
ent than the inclusive-OR operation; indeed, it is closer to the *exclusive-OR* operation.

The logical XOR (exclusive-or) operation is also a dyadic operation. It is defined as
follows:

\[
\begin{align*}
0 \text{ xor } 0 &= 0 \\
0 \text{ xor } 1 &= 1 \\
1 \text{ xor } 0 &= 1 \\
1 \text{ xor } 1 &= 0
\end{align*}
\]

The truth table for the XOR operation takes the following form:

**Table 4: XOR Truth Table**

<table>
<thead>
<tr>
<th>XOR</th>
<th>0</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
</tbody>
</table>
In English, the logical XOR operation is, “If the first operand or the second operand, but not both, is one, the result is one; otherwise the result is zero.” Note that the exclusive-or operation is closer to the English meaning of the word “or” than is the logical OR operation.

If one of the operands to the logical exclusive-OR operation is a one, the result is always the inverse of the other operand; that is, if one operand is one, the result is zero if the other operand is one and the result is one if the other operand is zero. If the first operand contains a zero, then the result is exactly the value of the second operand. This feature lets you selectively invert bits in a bit string.

The logical NOT operation is a monadic operation (meaning it accepts only one operand). It is:

\[
\text{NOT } 0 = 1 \\
\text{NOT } 1 = 0
\]

The truth table for the NOT operation takes the following form:

<table>
<thead>
<tr>
<th>NOT</th>
<th>0</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>1</td>
<td>0</td>
</tr>
</tbody>
</table>

### 1.6 Logical Operations on Binary Numbers and Bit Strings

As described in the previous section, the logical functions work only with single bit operands. Since the 80x86 uses groups of eight, sixteen, or thirty-two bits, we need to extend the definition of these functions to deal with more than two bits. Logical functions on the 80x86 operate on a bit-by-bit (or bitwise) basis. Given two values, these functions operate on bit zero producing bit zero of the result. They operate on bit one of the input values producing bit one of the result, etc. For example, if you want to compute the logical AND of the following two eight-bit numbers, you would perform the logical AND operation on each column independently of the others:

<table>
<thead>
<tr>
<th>1011</th>
<th>0101</th>
</tr>
</thead>
<tbody>
<tr>
<td>1110</td>
<td>1110</td>
</tr>
<tr>
<td>------</td>
<td>------</td>
</tr>
<tr>
<td>1010</td>
<td>0100</td>
</tr>
</tbody>
</table>

This bit-by-bit form of execution can be easily applied to the other logical operations as well.

Since we’ve defined logical operations in terms of binary values, you’ll find it much easier to perform logical operations on binary values than on values in other bases. Therefore, if you want to perform a logical operation on two hexadecimal numbers, you should convert them to binary first. This applies to most of the basic logical operations on binary numbers (e.g., AND, OR, XOR, etc.).

The ability to force bits to zero or one using the logical AND/OR operations and the ability to invert bits using the logical XOR operation is very important when working with strings of bits (e.g., binary numbers). These operations let you selectively manipulate certain bits within some value while leaving other bits unaffected. For example, if you have an eight-bit binary value ‘X’ and you want to guarantee that bits four through seven contain zeros, you could logically AND the value ‘X’ with the binary value 0000 1111. This

---

5. Monadic means the operator has one operand.
bitwise logical AND operation would force the H.O. four bits to zero and pass the L.O. four bits of ‘X’ through unchanged. Likewise, you could force the L.O. bit of ‘X’ to one and invert bit number two of ‘X’ by logically ORing ‘X’ with 0000 0001 and logically exclusive-ORing ‘X’ with 0000 0100, respectively. Using the logical AND, OR, and XOR operations to manipulate bit strings in this fashion is known as *masking* bit strings. We use the term *masking* because we can use certain values (one for AND, zero for OR/XOR) to ‘mask out’ certain bits from the operation when forcing bits to zero, one, or their inverse.

### 1.7 Signed and Unsigned Numbers

So far, we’ve treated binary numbers as unsigned values. The binary number ...00000 represents zero, ...00001 represents one, ...00010 represents two, and so on toward infinity. What about negative numbers? Signed values have been tossed around in previous sections and we’ve mentioned the two’s complement numbering system, but we haven’t discussed how to represent negative numbers using the binary numbering system. That is what this section is all about!

To represent signed numbers using the binary numbering system we have to place a restriction on our numbers: they must have a finite and fixed number of bits. As far as the 80x86 goes, this isn’t too much of a restriction, after all, the 80x86 can only address a finite number of bits. For our purposes, we’re going to severely limit the number of bits to eight, 16, 32, or some other small number of bits.

With a fixed number of bits we can only represent a certain number of objects. For example, with eight bits we can only represent 256 different objects. Negative values are objects in their own right, just like positive numbers. Therefore, we’ll have to use some of the 256 different values to represent negative numbers. In other words, we’ve got to use up some of the positive numbers to represent negative numbers. To make things fair, we’ll assign half of the possible combinations to the negative values and half to the positive values. So we can represent the negative values -128..-1 and the positive values 0..127 with a single eight bit byte. With a 16-bit word we can represent values in the range -32,768..+32,767. With a 32-bit double word we can represent values in the range -2,147,483,648..+2,147,483,647. In general, with \( n \) bits we can represent the signed values in the range \(-2^{n-1}\) to \(+2^{n-1}-1\).

Okay, so we can represent negative values. Exactly how do we do it? Well, there are many ways, but the 80x86 microprocessor uses the two’s complement notation. In the two’s complement system, the H.O. bit of a number is a *sign bit*. If the H.O. bit is zero, the number is positive; if the H.O. bit is one, the number is negative. Examples:

For 16-bit numbers:
- 8000h is negative because the H.O. bit is one.
- 100h is positive because the H.O. bit is zero.
- 7FFFh is positive.
- 0FFFh is negative.
- 0FFFh is positive.

If the H.O. bit is zero, then the number is positive and is stored as a standard binary value. If the H.O. bit is one, then the number is negative and is stored in the two’s complement form. To convert a positive number to its negative, two’s complement form, you use the following algorithm:

1) Invert all the bits in the number, i.e., apply the logical NOT function.

---

6. Technically, zero is neither positive nor negative. For technical reasons (due to the hardware involved), we’ll lump zero in with the positive numbers.
2) Add one to the inverted result.

For example, to compute the eight bit equivalent of -5:

```
0000 0101  Five (in binary).
1111 1010  Invert all the bits.
1111 1011  Add one to obtain result.
```

If we take minus five and perform the two’s complement operation on it, we get our original value, 00000101, back again, just as we expect:

```
1111 1011  Two’s complement for -5.
0000 0100  Invert all the bits.
0000 0101  Add one to obtain result (+5).
```

The following examples provide some positive and negative 16-bit signed values:

- **7FFFh**: +32767, the largest 16-bit positive number.
- **8000h**: -32768, the smallest 16-bit negative number.
- **4000h**: +16,384.

To convert the numbers above to their negative counterpart (i.e., to negate them), do the following:

**7FFFh**: 0111 1111 1111 1111 +32,767t
```
1000 0000 0000 0000  Invert all the bits (8000h)
1000 0000 0000 0001  Add one (8001h or -32,767t)
```

**8000h**: 1000 0000 0000 0000 -32,768t
```
0111 1111 1111 1111  Invert all the bits (7FFFh)
1000 0000 0000 0000  Add one (8000h or -32768t)
```

**4000h**: 0100 0000 0000 0000 16,384t
```
1011 1111 1111 1111  Invert all the bits (BFFFe)
1100 0000 0000 0000  Add one (0C000h or -16,384t)
```

8000h inverted becomes 7FFFh. After adding one we obtain 8000h! Wait, what’s going on here? -(-32,768) is -32,768? Of course not. But the value +32,768 cannot be represented with a 16-bit signed number, so we cannot negate the smallest negative value. If you attempt this operation, the 80x86 microprocessor will complain about signed arithmetic overflow.

Why bother with such a miserable numbering system? Why not use the H.O. bit as a sign flag, storing the positive equivalent of the number in the remaining bits? The answer lies in the hardware. As it turns out, negating values is the only tedious job. With the two’s complement system, most other operations are as easy as the binary system. For example, suppose you were to perform the addition 5+(-5). The result is zero. Consider what happens when we add these two values in the two’s complement system:

```
00000101
11111011
--------
1 00000000
```

We end up with a carry into the ninth bit and all other bits are zero. As it turns out, if we ignore the carry out of the H.O. bit, adding two signed values always produces the correct result when using the two’s complement numbering system. This means we can use the same hardware for signed and unsigned addition and subtraction. This wouldn’t be the case with some other numbering systems.

Except for the questions at the end of this chapter, you will not need to perform the two’s complement operation by hand. The 80x86 microprocessor provides an instruction, NEG (negate), which performs this operation for you. Furthermore, all the hexadecimal
calculators will perform this operation by pressing the change sign key (+/- or CHS). Nevertheless, performing a two’s complement by hand is easy, and you should know how to do it.

Once again, you should note that the data represented by a set of binary bits depends entirely on the context. The eight bit binary value 11000000b could represent an IBM/ASCII character, it could represent the unsigned decimal value 192, or it could represent the signed decimal value -64, etc. As the programmer, it is your responsibility to use this data consistently.

1.8 Sign and Zero Extension

Since two’s complement format integers have a fixed length, a small problem develops. What happens if you need to convert an eight bit two’s complement value to 16 bits? This problem, and its converse (converting a 16 bit value to eight bits) can be accomplished via sign extension and contraction operations. Likewise, the 80x86 works with fixed length values, even when processing unsigned binary numbers. Zero extension lets you convert small unsigned values to larger unsigned values.

Consider the value “-64”. The eight bit two’s complement value for this number is 0C0h. The 16-bit equivalent of this number is 0FFC0h. Now consider the value “+64”. The eight and 16 bit versions of this value are 040h and 0040h. The difference between the eight and 16 bit numbers can be described by the rule: “If the number is negative, the H.O. byte of the 16 bit number contains 0FFh; if the number is positive, the H.O. byte of the 16 bit quantity is zero.”

To sign extend a value from some number of bits to a greater number of bits is easy, just copy the sign bit into all the additional bits in the new format. For example, to sign extend an eight bit number to a 16 bit number, simply copy bit seven of the eight bit number into bits 8..15 of the 16 bit number. To sign extend a 16 bit number to a double word, simply copy bit 15 into bits 16..31 of the double word.

Sign extension is required when manipulating signed values of varying lengths. Often you’ll need to add a byte quantity to a word quantity. You must sign extend the byte quantity to a word before the operation takes place. Other operations (multiplication and division, in particular) may require a sign extension to 32-bits. You must not sign extend unsigned values.

Examples of sign extension:

<table>
<thead>
<tr>
<th>Eight Bits</th>
<th>Sixteen Bits</th>
<th>Thirty-two Bits</th>
</tr>
</thead>
<tbody>
<tr>
<td>80h</td>
<td>FF80h</td>
<td>FFFFFFF80h</td>
</tr>
<tr>
<td>28h</td>
<td>0028h</td>
<td>00000028h</td>
</tr>
<tr>
<td>9Ah</td>
<td>FF9Ah</td>
<td>FFFFFF9Ah</td>
</tr>
<tr>
<td>7Fh</td>
<td>007Fh</td>
<td>0000007Fh</td>
</tr>
<tr>
<td>---</td>
<td>1020h</td>
<td>00001020h</td>
</tr>
<tr>
<td>---</td>
<td>8088h</td>
<td>FFFF8088h</td>
</tr>
</tbody>
</table>

To extend an unsigned byte you must zero extend the value. Zero extension is very easy – just store a zero into the H.O. byte(s) of the smaller operand. For example, to zero extend the value 82h to 16-bits you simply add a zero to the H.O. byte yielding 0082h.

<table>
<thead>
<tr>
<th>Eight Bits</th>
<th>Sixteen Bits</th>
<th>Thirty-two Bits</th>
</tr>
</thead>
<tbody>
<tr>
<td>80h</td>
<td>0080h</td>
<td>00000080h</td>
</tr>
<tr>
<td>28h</td>
<td>0028h</td>
<td>00000028h</td>
</tr>
<tr>
<td>9Ah</td>
<td>009Ah</td>
<td>0000009Ah</td>
</tr>
<tr>
<td>7Fh</td>
<td>007Fh</td>
<td>0000007Fh</td>
</tr>
<tr>
<td>---</td>
<td>1020h</td>
<td>00001020h</td>
</tr>
<tr>
<td>---</td>
<td>8088h</td>
<td>00008088h</td>
</tr>
</tbody>
</table>

Sign contraction, converting a value with some number of bits to the identical value with a fewer number of bits, is a little more troublesome. Sign extension never fails. Given an m-bit signed value you can always convert it to an n-bit number (where n > m) using
sign extension. Unfortunately, given an \( n \)-bit number, you cannot always convert it to an \( m \)-bit number if \( m < n \). For example, consider the value -448. As a 16-bit hexadecimal number, its representation is 0FE40h. Unfortunately, the magnitude of this number is too great to fit into an eight bit value, so you cannot sign contract it to eight bits. This is an example of an overflow condition that occurs upon conversion.

To properly sign contract one value to another, you must look at the H.O. byte(s) that you want to discard. The H.O. bytes you wish to remove must all contain either zero or 0FFh. If you encounter any other values, you cannot contract it without overflow. Finally, the H.O. bit of your resulting value must match every bit you’ve removed from the number. Examples (16 bits to eight bits):

- FF80h can be sign contracted to 80h
- 0040h can be sign contracted to 40h
- FE40h cannot be sign contracted to 8 bits.
- 0100h cannot be sign contracted to 8 bits.

### 1.9 Shifts and Rotates

Another set of logical operations which apply to bit strings are the **shift** and **rotate** operations. These two categories can be further broken down into left **shifts**, left **rotates**, right **shifts**, and right **rotates**. These operations turn out to be extremely useful to assembly language programmers.

The left shift operation moves each bit in a bit string one position to the left (see Figure 1.8).

![Figure 1.8: Shift Left Operation](image)

Bit zero moves into bit position one, the previous value in bit position one moves into bit position two, etc. There are, of course, two questions that naturally arise: “What goes into bit zero?” and “Where does bit seven wind up?” Well, that depends on the context. We’ll shift the value zero into the L.O. bit, and the previous value of bit seven will be the **carry** out of this operation.

Note that shifting a value to the left is the same thing as multiplying it by its radix. For example, shifting a decimal number one position to the left (adding a zero to the right of the number) effectively multiplies it by ten (the radix):

\[
1234 \text{ SHL 1} = 12340 \quad (\text{SHL 1 = shift left one position})
\]

Since the radix of a binary number is two, shifting it left multiplies it by two. If you shift a binary value to the left twice, you multiply it by two twice (i.e., you multiply it by four). If you shift a binary value to the left three times, you multiply it by eight (\(2^3\)). In general, if you shift a value to the left \( n \) times, you multiply that value by \( 2^n \).

A right shift operation works the same way, except we’re moving the data in the opposite direction. Bit seven moves into bit six, bit six moves into bit five, bit five moves into bit four, etc. During a right shift, we’ll move a zero into bit seven, and bit zero will be the **carry** out of the operation (see Figure 1.9).

Since a left shift is equivalent to a multiplication by two, it should come as no surprise that a right shift is roughly comparable to a division by two (or, in general, a division by the radix of the number). If you perform \( n \) right shifts, you will divide that number by \( 2^n \).
There is one problem with shift rights with respect to division: as described above a shift right is only equivalent to an unsigned division by two. For example, if you shift the unsigned representation of 254 (0FEh) one place to the right, you get 127 (07Fh), exactly what you would expect. However, if you shift the binary representation of -2 (0FEh) to the right one position, you get 127 (07Fh), which is not correct. This problem occurs because we’re shifting a zero into bit seven. If bit seven previously contained a one, we’re changing it from a negative to a positive number. Not a good thing when dividing by two.

To use the shift right as a division operator, we must define a third shift operation: arithmetic shift right. An arithmetic shift right works just like the normal shift right operation (a logical shift right) with one exception: instead of shifting a zero into bit seven, an arithmetic shift right operation leaves bit seven alone, that is, during the shift operation it does not modify the value of bit seven as Figure 1.10 shows.

This generally produces the result you expect. For example, if you perform the arithmetic shift right operation on -2 (0FEh) you get -1 (0FFh). Keep one thing in mind about arithmetic shift right, however. This operation always rounds the numbers to the closest integer which is less than or equal to the actual result. Based on experiences with high level programming languages and the standard rules of integer truncation, most people assume this means that a division always truncates towards zero. But this simply isn’t the case. For example, if you apply the arithmetic shift right operation on -1 (0FFh), the result is -1, not zero. -1 is less than zero so the arithmetic shift right operation rounds towards minus one. This is not a “bug” in the arithmetic shift right operation. This is the way integer division typically gets defined. The 80x86 integer division instruction also produces this result.

Another pair of useful operations are rotate left and rotate right. These operations behave like the shift left and shift right operations with one major difference: the bit shifted out from one end is shifted back in at the other end.

---

7. There is no need for an arithmetic shift left. The standard shift left operation works for both signed and unsigned numbers, assuming no overflow occurs.
1.10 Bit Fields and Packed Data

Although the 80x86 operates most efficiently on byte, word, and double word data types, occasionally you'll need to work with a data type that uses some number of bits other than eight, 16, or 32. For example, consider a date of the form “4/2/88”. It takes three numeric values to represent this date: a month, day, and year value. Months, of course, take on the values 1..12. It will require at least four bits (maximum of sixteen different values) to represent the month. Days range between 1..31. So it will take five bits (maximum of 32 different values) to represent the day entry. The year value, assuming that we’re working with values in the range 0..99, requires seven bits (which can be used to represent up to 128 different values). Four plus five plus seven is 16 bits, or two bytes. In other words, we can pack our date data into two bytes rather than the three that would be required if we used a separate byte for each of the month, day, and year values. This saves one byte of memory for each date stored, which could be a substantial saving if you need to store a lot of dates. The bits could be arranged as shown in .

![Figure 1.12: Rotate Right Operation](image)

![Figure 1.13: Packed Date Format](image)

MMMM represents the four bits making up the month value, DDDDD represents the five bits making up the day, and YYYYYYYY is the seven bits comprising the year. Each collection of bits representing a data item is a *bit field*. April 2nd, 1988 would be represented as 4158h:

```
0100 00010 1011000 = 0100 0001 0101 1000b or 4158h
```

Although packed values are *space efficient* (that is, very efficient in terms of memory usage), they are computationally *inefficient* (slow!). The reason? It takes extra instructions to unpack the data packed into the various bit fields. These extra instructions take additional time to execute (and additional bytes to hold the instructions); hence, you must carefully consider whether packed data fields will save you anything.

Examples of practical packed data types abound. You could pack eight boolean values into a single byte, you could pack two BCD digits into a byte, etc.

1.11 The ASCII Character Set

The ASCII character set (excluding the extended characters defined by IBM) is divided into four groups of 32 characters. The first 32 characters, ASCII codes 0 through
1Fh (31), form a special set of non-printing characters called the control characters. We call them control characters because they perform various printer/display control operations rather than displaying symbols. Examples include carriage return, which positions the cursor to the left side of the current line of characters, line feed (which moves the cursor down one line on the output device), and back space (which moves the cursor back one position to the left). Unfortunately, different control characters perform different operations on different output devices. There is very little standardization among output devices. To find out exactly how a control character affects a particular device, you will need to consult its manual.

The second group of 32 ASCII character codes comprise various punctuation symbols, special characters, and the numeric digits. The most notable characters in this group include the space character (ASCII code 20h) and the numeric digits (ASCII codes 30h..39h). Note that the numeric digits differ from their numeric values only in the H.O. nibble. By subtracting 30h from the ASCII code for any particular digit you can obtain the numeric equivalent of that digit.

The third group of 32 ASCII characters is reserved for the upper case alphabetic characters. The ASCII codes for the characters “A”..”Z” lie in the range 41h..5Ah (65..90). Since there are only 26 different alphabetic characters, the remaining six codes hold various special symbols.

The fourth, and final, group of 32 ASCII character codes are reserved for the lower case alphabetic symbols, five additional special symbols, and another control character (delete). Note that the lower case character symbols use the ASCII codes 61h..7Ah. If you convert the codes for the upper and lower case characters to binary, you will notice that the upper case symbols differ from their lower case equivalents in exactly one bit position. For example, consider the character code for “E” and “e” in Figure 1.14.

![Figure 1.14: ASCII Codes for “E” and “e”](image)

The only place these two codes differ is in bit five. Upper case characters always contain a zero in bit five; lower case alphabetic characters always contain a one in bit five. You can use this fact to quickly convert between upper and lower case. If you have an upper case character you can force it to lower case by setting bit five to one. If you have a lower case character and you wish to force it to upper case, you can do so by setting bit five to zero. You can toggle an alphabetic character between upper and lower case by simply inverting bit five.

Indeed, bits five and six determine which of the four groups in the ASCII character set you’re in:

---

8. Historically, carriage return refers to the *paper carriage* used on typewriters. A carriage return consisted of physically moving the carriage all the way to the right so that the next character typed would appear at the left hand side of the paper.
So you could, for instance, convert any upper or lower case (or corresponding special) character to its equivalent control character by setting bits five and six to zero.

Consider, for a moment, the ASCII codes of the numeric digit characters:

<table>
<thead>
<tr>
<th>Char</th>
<th>Dec</th>
<th>Hex</th>
</tr>
</thead>
<tbody>
<tr>
<td>“0”</td>
<td>48</td>
<td>30h</td>
</tr>
<tr>
<td>“1”</td>
<td>49</td>
<td>31h</td>
</tr>
<tr>
<td>“2”</td>
<td>50</td>
<td>32h</td>
</tr>
<tr>
<td>“3”</td>
<td>51</td>
<td>33h</td>
</tr>
<tr>
<td>“4”</td>
<td>52</td>
<td>34h</td>
</tr>
<tr>
<td>“5”</td>
<td>53</td>
<td>35h</td>
</tr>
<tr>
<td>“6”</td>
<td>54</td>
<td>36h</td>
</tr>
<tr>
<td>“7”</td>
<td>55</td>
<td>37h</td>
</tr>
<tr>
<td>“8”</td>
<td>56</td>
<td>38h</td>
</tr>
<tr>
<td>“9”</td>
<td>57</td>
<td>39h</td>
</tr>
</tbody>
</table>

The decimal representations of these ASCII codes are not very enlightening. However, the hexadecimal representation of these ASCII codes reveals something very important – the L.O. nibble of the ASCII code is the binary equivalent of the represented number. By stripping away (i.e., setting to zero) the H.O. nibble of a numeric character, you can convert that character code to the corresponding binary representation. Conversely, you can convert a binary value in the range 0..9 to its ASCII character representation by simply setting the H.O. nibble to three. Note that you can use the logical-AND operation to force the H.O. bits to zero; likewise, you can use the logical-OR operation to force the H.O. bits to 0011 (three).

Note that you cannot convert a string of numeric characters to their equivalent binary representation by simply stripping the H.O. nibble from each digit in the string. Converting 123 (31h 32h 33h) in this fashion yields three bytes: 010203h, not the correct value which is 7Bh. Converting a string of digits to an integer requires more sophistication than this; the conversion above works only for single digits.

Bit seven in standard ASCII is always zero. This means that the ASCII character set consumes only half of the possible character codes in an eight bit byte. IBM uses the remaining 128 character codes for various special characters including international characters (those with accents, etc.), math symbols, and line drawing characters. Note that these extra characters are a non-standard extension to the ASCII character set. Of course, the name IBM has considerable clout, so almost all modern personal computers based on the 80x86 with a video display support the extended IBM/ASCII character set. Most printers support IBM’s character set as well.
Should you need to exchange data with other machines which are not PC-compatible, you have only two alternatives: stick to standard ASCII or ensure that the target machine supports the extended IBM-PC character set. Some machines, like the Apple Macintosh, do not provide native support for the extended IBM-PC character set; however you may obtain a PC font which lets you display the extended character set. Other machines (e.g., Amiga and Atari ST) have similar capabilities. However, the 128 characters in the standard ASCII character set are the only ones you should count on transferring from system to system.

Despite the fact that it is a "standard", simply encoding your data using standard ASCII characters does not guarantee compatibility across systems. While it's true that an "A" on one machine is most likely an "A" on another machine, there is very little standardization across machines with respect to the use of the control characters. Indeed, of the 32 control codes plus delete, there are only four control codes commonly supported – backspace (BS), tab, carriage return (CR), and line feed (LF). Worse still, different machines often use these control codes in different ways. End of line is a particularly troublesome example. MS-DOS, CP/M, and other systems mark end of line by the two-character sequence CR/LF. Apple Macintosh, Apple II, and many other systems mark the end of line by a single CR character. UNIX systems mark the end of a line with a single LF character. Needless to say, attempting to exchange simple text files between such systems can be an experience in frustration. Even if you use standard ASCII characters in all your files on these systems, you will still need to convert the data when exchanging files between them. Fortunately, such conversions are rather simple.

Despite some major shortcomings, ASCII data is the standard for data interchange across computer systems and programs. Most programs can accept ASCII data; likewise most programs can produce ASCII data. Since you will be dealing with ASCII characters in assembly language, it would be wise to study the layout of the character set and memorize a few key ASCII codes (e.g., "0", "A", "a", etc.).

1.12 Summary

Most modern computer systems use the binary numbering system to represent values. Since binary values are somewhat unwieldy, we'll often use the hexadecimal representation for those values. This is because it is very easy to convert between hexadecimal and binary, unlike the conversion between the more familiar decimal and binary systems. A single hexadecimal digit consumes four binary digits (bits), and we call a group of four bits a nibble. See:

- "The Binary Numbering System" on page 12
- "Binary Formats" on page 13
- "The Hexadecimal Numbering System" on page 17

The 80x86 works best with groups of bits which are eight, 16, or 32 bits long. We call objects of these sizes bytes, words, and double words, respectively. With a byte, we can represent any one of 256 unique values. With a word we can represent one of 65,536 different values. With a double word we can represent over four billion different values. Often we simply represent integer values (signed or unsigned) with bytes, words, and double words; however we'll often represent other quantities as well. See:

- "Data Organization" on page 13
- "Bytes" on page 14
- "Words" on page 15
- "Double Words" on page 16

In order to talk about specific bits within a nibble, byte, word, double word, or other structure, we'll number the bits starting at zero (for the least significant bit) on up to n-1
(where \( n \) is the number of bits in the object). We’ll also number nibbles, bytes, and words in large structures in a similar fashion. See:

- “Binary Formats” on page 13

There are many operations we can perform on binary values including normal arithmetic (\(+, -, \times, \text{ and } /\)) and the logical operations (AND, OR, XOR, NOT, Shift Left, Shift Right, Rotate Left, and Rotate Right). Logical AND, OR, XOR, and NOT are typically defined for single bit operations. We can extend these to \( n \) bits by performing bitwise operations. The shifts and rotates are always defined for a fixed length string of bits. See:

- “Arithmetic Operations on Binary and Hexadecimal Numbers” on page 19
- “Logical Operations on Bits” on page 20
- “Logical Operations on Binary Numbers and Bit Strings” on page 22
- “Shifts and Rotates” on page 26

There are two types of integer values which we can represent with binary strings on the 80x86: unsigned integers and signed integers. The 80x86 represents unsigned integers using the standard binary format. It represents signed integers using the two’s complement format. While unsigned integers may be of arbitrary length, it only makes sense to talk about fixed length signed binary values. See:

- “Signed and Unsigned Numbers” on page 23
- “Sign and Zero Extension” on page 25

Often it may not be particularly practical to store data in groups of eight, 16, or 32 bits. To conserve space you may want to pack various pieces of data into the same byte, word, or double word. This reduces storage requirements at the expense of having to perform extra operations to pack and unpack the data. See:

- “Bit Fields and Packed Data” on page 28

Character data is probably the most common data type encountered besides integer values. The IBM PC and compatibles use a variant of the ASCII character set – the extended IBM/ASCII character set. The first 128 of these characters are the standard ASCII characters, 128 are special characters created by IBM for international languages, mathematics, and line drawing. Since the use of the ASCII character set is so common in modern programs, familiarity with this character set is essential. See:

- “The ASCII Character Set” on page 28
1.13 Laboratory Exercises

Accompanying this text is a significant amount of software. This software is divided into four basic categories: source code for examples appearing throughout this text, the UCR Standard Library for 80x86 assembly language programmers, sample code you modify for various laboratory exercises, and application software to support various laboratory exercises. This software has been written using assembly language, C++, Flex/Bison, and Delphi (object Pascal). Most of the application programs include source code as well as executable code.

Much of the software accompanying this text runs under Windows 3.1, Windows 95, or Windows NT. Some software, however, directly manipulates the hardware and will only run under DOS or a DOS box in Windows 3.1. This text assumes that you are familiar with the DOS and Windows operating systems; if you are unfamiliar with DOS or Windows operation, you should refer to an appropriate text on those systems for additional details.

1.13.1 Installing the Software

The software accompanying this text is generally supplied on CD-ROM. You can use most of it as-is directly off the CD-ROM. However, for speed and convenience you will probably want to install the software on a hard disk. To do this, you will need to create two subdirectories in the root directory on your hard drive: ARTOFASM and STDLIB. The ARTOFASM directory will contain the files specific to this text book, the STDLIB directory will contain the files associated with the UCR Standard Library for 80x86 assembly language programmers. Once you create these two subdirectories, copy all the files and subdirectories from the corresponding directories on the CD to your hard disk. From DOS (or a DOS window), you can use the following XCOPY commands to accomplish this:

```
xcopy r:\artofasm\*.*   c:\artofasm   /s
xcopy r:\stdlib\*.*   c:\stdlib   /s
```

These commands assume that your CD-ROM is drive R: and you are installing the software on the C: hard disk. They also assume that you have created the ARTOFASM and STDLIB subdirectories prior to executing the XCOPY commands.

To use the Standard Library in programming projects, you will need to add or modify two lines in your AUTOEXEC.BAT file. If similar lines are not already present, add the following two lines to your AUTOEXEC.BAT file:

```
set lib=c:\stdlib\lib
set include=c:\stdlib\include
```

These commands tell MASM (the Microsoft Macro Assembler) where it can find the library and include files for the UCR Standard Library. Without these lines, MASM will report an error anytime you use the standard library routines in your programs.

If there are already a “set include = ...” and “set lib=...” lines in your AUTOEXEC.BAT file, you should not replace them with the lines above. Instead, you should append the string “c:\stdlib\lib” to the end of the existing “set lib=...” statement and “c:\stdlib\include” to the end of the existing “set include=...” statement. Several languages (like C++) also use these “set” statements; if you arbitrarily replace them with the statements above, your assembly language programs will work fine, but any attempt to compile a C++ (or other language) program may fail.

---

9. It is also available via anonymous ftp, although there are many files associated with this text.
10. If you are using this software in a laboratory at school, your instructor has probably installed this software on the machines in the laboratory. As a general rule, you should never install software on machines in the laboratory. Check with your laboratory instruction before installing this software on machines in the laboratory.
If you forget to put these lines in your AUTOEXEC.BAT file, you can temporarily (until the next time you boot the system) issue these commands by simply typing them at the DOS command line prompt. By typing “set” by itself on the command line prompt, you can see if these set commands are currently active.

If you do not have a CD-ROM player, you can obtain the software associated with this textbook via anonymous ftp from cs.ucr.edu. Check in the “/pub/pc/ibmpc” subdirectory. The files on the ftp server will be compressed. A “README” file will describe how to decompress the data.

The STDLIB directory you’ve created holds the source and library files for the UCR Standard Library for 80x86 Assembly Language Programmers. This is a core set of assembly language subroutines you can call that mimic many of the routines in the C standard library. These routines greatly simplify writing programs in assembly language. Furthermore, they are public domain so you can use them in any programs you write without fear of licensing restrictions.

The ARTOFASM directory contains files specific to this text. Within the ARTOFASM directory you will see a sequence of subdirectories named ch1, ch2, ch3, etc. These subdirectories contain the files associated with Chapter One, Chapter Two, and so on. Within some of these subdirectories, you will find two subdirectories named “DOS” and “WINDOWS”. If these subdirectories are present, they separate those files that must run under MS-Windows from those that run under DOS. Many of the DOS programs require a “real-mode” environment and will not run in a DOS box window in Windows 95 or Windows NT. You will need to run this software directory from MS-DOS. The Windows applications require a color monitor.

There is often a third subdirectory present in each chapter directory: SOURCES. This subdirectory contains the source listings (where appropriate or feasible) to the software for that chapter. Most of the software for this text is written in assembly language using MASM 6.x, generic C++, Turbo Pascal, or Borland Delphi (visual object Pascal). If you are interested in seeing how the software operates, you can look in this subdirectory.

This text assumes you already know how to run programs from MS-DOS and Windows and you are familiar with common DOS and Windows terminology. It also assumes you know some simple MS-DOS commands like DIR, COPY, DEL, RENAME, and so on. If you are new to Windows and DOS, you should pick up an appropriate reference manual on these operating systems.

The files for Chapter One’s laboratory exercises appear in the ARTOFASM\CH1 subdirectory. These are all Windows programs, so you will need to be running Windows 3.1, Windows 95, Windows NT, or some later (and compatible) version of Windows to run these programs.

1.13.2 Data Conversion Exercises

In this exercise you will be using the “convert.exe” program found in the ARTOFASM\CH1 subdirectory. This program displays and converts 16-bit integers using signed decimal, unsigned decimal, hexadecimal, and binary notation.

When you run this program it opens a window with four edit boxes. (one for each data type). Changing a value in one of the edit boxes immediately updates the values in the other boxes so they all display their corresponding representations for the new value. If you make a mistake on data entry, the program beeps and turns the edit box red until you correct the mistake. Note that you can use the mouse, cursor control keys, and the editing keys (e.g., DEL and Backspace) to change individual values in the edit boxes.

For this exercise and your laboratory report, you should explore the relationship between various binary, hexadecimal, unsigned decimal, and signed decimal values. For example, you should enter the unsigned decimal values 1, 2, 4, 8, 16, 32, 64, 128, 256, 512, 1024, 2048, 4096, 8192, 16384, and 32768 and comment on the values that appear in the in the other text boxes.
The primary purpose of this exercise is to familiarize yourself with the decimal equivalents of some common binary and hexadecimal values. In your lab report, for example, you should explain what is special about the binary (and hexadecimal) equivalents of the decimal numbers above.

Another set of experiments to try is to choose various binary numbers that have exactly two bits set, e.g., 11, 110, 1100, 1 1000, 11 0000, etc. Be sure to comment on the decimal and hexadecimal results these inputs produce.

Try entering several binary numbers where the L.O. eight bits are all zero. Comment on the results in your lab report. Try the same experiment with hexadecimal numbers using zeros for the L.O. digit or the two L.O. digits.

You should also experiment with negative numbers in the signed decimal text entry box; try using values like -1, -2, -3, -256, -1024, etc. Explain the results you obtain using your knowledge of the two’s complement numbering system.

Try entering even and odd numbers in unsigned decimal. Discover and describe the difference between even and odd numbers in their binary representation. Try entering multiples of other values (e.g., for three: 3, 6, 9, 12, 15, 18, 21, ...) and see if you can detect a pattern in the binary results.

Verify the hexadecimal <> binary conversion this chapter describes. In particular, enter the same hexadecimal digit in each of the four positions of a 16-bit value and comment on the position of the corresponding bits in the binary representation. Try several entering binary values like 1111, 11110, 111100, 1111000, and 11110000. Explain the results you get and describe why you should always extend binary values so their length is an even multiple of four before converting them.

In your lab report, list the experiments above plus several you devise yourself. Explain the results you expect and include the actual results that the convert.exe program produces. Explain any insights you have while using the convert.exe program.

### 1.13.3 Logical Operations Exercises

The logical.exe program is a simple calculator that computes various logical functions. It allows you to enter binary or hexadecimal values and then it computes the result of some logical operation on the inputs. The calculator supports the dyadic logical AND, OR, and XOR. It also supports the monadic NOT, NEG (two’s complement), SHL (shift left), SHR (shift right), ROL (rotate left), and ROR (rotate right).

When you run the logical.exe program it displays a set of buttons on the left hand side of the window. These buttons let you select the calculation. For example, pressing the AND button instructs the calculator to compute the logical AND operation between the two input values. If you select a monadic (unary) operation like NOT, SHL, etc., then you may only enter a single value; for the dyadic operations, both sets of text entry boxes will be active.

The logical.exe program lets you enter values in binary or hexadecimal. Note that this program automatically converts any changes in the binary text entry window to hexadecimal and updates the value in the hex entry edit box. Likewise, any changes in the hexadecimal text entry box are immediately reflected in the binary text box. If you enter an illegal value in a text entry box, the logical.exe program will turn the box red until you correct the problem.

For this laboratory exercise, you should explore each of the bitwise logical operations. Create several experiments by carefully choosing some values, manually compute the result you expect, and then run the experiment using the logical.exe program to verify your results. You should especially experiment with the masking capabilities of the logical AND, OR, and XOR operations. Try logically ANDing, ORing, and XORing different values with values like 000F, 00FF, 00F0, 0FFF, FF00, etc. Report the results and comment on them in your laboratory report.
Some experiments you might want to try, in addition to those you devise yourself, include the following:

- Devise a mask to convert ASCII values ‘0’..'9' to their binary integer counterparts using the logical AND operation. Try entering the ASCII codes of each of these digits when using this mask. Describe your results. What happens if you enter non-digit ASCII codes?
- Devise a mask to convert integer values in the range 0..9 to their corresponding ASCII codes using the logical OR operation. Enter each of the binary values in the range 0..9 and describe your results. What happens if you enter values outside the range 0..9? In particular, what happens if you enter values outside the range 0h..0fh?
- Devise a mask to determine whether a 16-bit integer value is positive or negative using the logical AND operation. The result should be zero if the number is positive (or zero) and it should be non-zero if the number is negative. Enter several positive and negative values to test your mask. Explain how you could use the AND operation to test any single bit to determine if it is zero or one.
- Devise a mask to use with the logical XOR operation that will produce the same result on the second operand as applying the logical NOT operator to that second operand.
- Verify that the SHL and SHR operators correspond to an integer multiplication by two and an integer division by two, respectively. What happens if you shift data out of the H.O. or L.O. bits? What does this correspond to in terms of integer multiplication and division?
- Apply the ROL operation to a set of positive and negative numbers. Based on your observations in Section 1.13.3, what can you say will about the result when you rotate left a negative number or a positive number?
- Apply the NEG and NOT operators to a value. Discuss the similarity and the difference in their results. Describe this difference based on your knowledge of the two’s complement numbering system.

### 1.13.4 Sign and Zero Extension Exercises

The signext.exe program accepts eight-bit binary or hexadecimal values then sign and zero extends them to 16 bits. Like the logical.exe program, this program lets you enter a value in either binary or hexadecimal and immediate zero and sign extends that value.

For your laboratory report, provide several eight-bit input values and describe the results you expect. Run these values through the signext.exe program and verify the results. For each experiment you run, be sure to list all the results in your lab report. Be sure to try values like 0, 7fh, 80h, and 0ffh.

While running these experiments, discover which hexadecimal digits appearing in the H.O. nibble produce negative 16-bit numbers and which produce positive 16-bit values. Document this set in your lab report.

Enter sets of values like (1,10), (2,20), (3,30), ..., (7,70), (8,80), (9,90), (A,A0), ..., (F,F0). Explain the results you get in your lab report. Why does “F” sign extend with zeros while “F0” sign extends with ones?

Explain in your lab report how one would sign or zero extend 16 bit values to 32 bit values. Explain why zero extension or sign extension is useful.
1.13.5 Packed Data Exercises

The packdata.exe program uses the Date data type appearing in this chapter (see “Bit Fields and Packed Data” on page 28). It lets you input a date value in binary or decimal and it packs that date into a single 16-bit value.

When you run this program, it will give you a window with six data entry boxes: three to enter the date in decimal form (month, day, year) and three text entry boxes that let you enter the date in binary form. The month value should be in the range 1..12, the day value should be in the range 1..31, and the year value should be in the range 0..99. If you enter a value outside this range (or some other illegal value), then the packdata.exe program will turn the data entry box red until you correct the problem.

Choose several dates for your experiments and convert these dates to the 16-bit packed binary form by hand (if you have trouble with the decimal to binary conversion, use the conversion program from the first set of exercises in this laboratory). Then run these dates through the packdata.exe program to verify your answer. Be sure to include all program output in your lab report.

At a bare minimum, you should include the following dates in your experiments:

2/4/68, 1/1/80, 8/16/64, 7/20/60, 11/2/72, 12/25/99, Today’s Date, a birthday (not necessarily yours), the due date on your lab report.
1.14 Questions

1) Convert the following decimal values to binary:

   a) 128   b) 4096   c) 256   d) 65536   e) 254
   f) 9     g) 1024   h) 15    i) 344    j) 998
   k) 255   l) 512    m) 1023   n) 2048   o) 4095
   p) 8192  q) 16,384 r) 32,768 s) 6,334   t) 12,334
   u) 23,465 v) 5,643  w) 464    x) 67     y) 888

2) Convert the following binary values to decimal:

   a) 1001 1001   b) 1001 1101   c) 1100 0011   d) 0000 1001   e) 1111 1111
   f) 0000 1111   g) 0111 1111   h) 1010 0101   i) 0100 0101   j) 0101 1010
   k) 1111 0000   l) 1011 1101   m) 1100 0010   n) 0111 1110   o) 1110 1111
   p) 0001 1000   q) 1001 1111  r) 0100 0010   s) 1101 1100   t) 1111 0001
   u) 0110 1001   v) 0101 1011   w) 1011 1001   x) 1110 0110   y) 1001 0111

3) Convert the binary values in problem 2 to hexadecimal.

4) Convert the following hexadecimal values to binary:

   a) 0ABCD   b) 1024  c) 0DEAD   d) 0ADD    e) 0BEEF
   f) 8       g) 05AAF h) 0FFF i) 0ACDB j) 0CDBA
   k) 0FEBA   l) 35   m) 0BA  n) 0ABA o) 0BAD
   p) 0DAB   q) 4321 r) 334  s) 45  t) 0E65
   u) 0BEAD   v) 0ABE w) 0DEAF x) 0DAD y) 9876

Perform the following hex computations (leave the result in hex):

5) 1234 +9876
6) 0FFF - 0F34
7) 100 - 1
8) 0FFE - 1

9) What is the importance of a nibble?

10) How many hexadecimal digits in:

    a) a byte   b) a word   c) a double word

11) How many bits in a:

    a) nibble   b) byte    c) word   d) double word

12) Which bit (number) is the H.O. bit in a:

    a) nibble   b) byte    c) word   d) double word

13) What character do we use as a suffix for hexadecimal numbers? Binary numbers? Decimal numbers?

14) Assuming a 16-bit two’s complement format, determine which of the values in question 4 are positive and which are negative.

15) Sign extend all of the values in question two to sixteen bits. Provide your answer in hex.
16) Perform the bitwise AND operation on the following pairs of hexadecimal values. Present your answer in hex. (Hint: convert hex values to binary, do the operation, then convert back to hex).

   a) 0FF00, 0FF0   b) 0F00F, 1234   c) 4321, 1234   d) 2341, 3241   e) 0FFFF, 0EDCB
   f) 1111, 5789   g) 0FABA, 4322   h) 5523, 0F572   i) 2355, 7466   j) 4765, 6543
   k) 0ABCD, 0FEDC   l) 0DDDD, 1234   m) 0CCCC, 0ABCDA   n) 0BBBB, 1234   o) 0AAAA, 1234
   p) 0EEEE, 1248   q) 8888, 1248   r) 8086, 124F   s) 8086, 0CFA7   t) 8765, 3456
   u) 7089, 0FEDC   v) 2435, 0BCDE   w) 6355, 0EFDC   x) 0CBA, 6884   y) 0AC7, 365

17) Perform the logical OR operation on the above pairs of numbers.

18) Perform the logical XOR operation on the above pairs of numbers.

19) Perform the logical NOT operation on all the values in question four. Assume all values are 16 bits.

20) Perform the two’s complement operation on all the values in question four. Assume 16 bit values.

21) Sign extend the following hexadecimal values from eight to sixteen bits. Present your answer in hex.

   a) FF   b) 82   c) 12   d) 56   e) 98
   f) BF   g) 0F   h) 78   i) 7F   j) F7
   k) 0E   l) AE   m) 45   n) 93   o) C0
   p) 8F   q) DA   r) 1D   s) 0D   t) DE
   u) 54   v) 45   w) F0   x) AD   y) DD

22) Sign contract the following values from sixteen bits to eight bits. If you cannot perform the operation, explain why.

   a) FF00   b) FF12   c) FFF0   d) 12   e) 80
   f) FFFF   g) FF88   h) FF7F   i) 7F   j) 2
   k) 8080   l) 80FF   m) FF80   n) FF   o) 8
   p) F   q) 1   r) 834   s) 34   t) 23
   u) 67   v) 89   w) 98   x) FF98   y) F98

23) Sign extend the 16-bit values in question 22 to 32 bits.

24) Assuming the values in question 22 are 16-bit values, perform the left shift operation on them.

25) Assuming the values in question 22 are 16-bit values, perform the right shift operation on them.

26) Assuming the values in question 22 are 16-bit values, perform the rotate left operation on them.

27) Assuming the values in question 22 are 16-bit values, perform the rotate right operation on them.

28) Convert the following dates to the packed format described in this chapter (see “Bit Fields and Packed Data” on page 28). Present your values as a 16-bit hex number.

   a) 1/1/92   b) 2/4/56   c) 6/19/60   d) 6/16/86   e) 1/1/99

29) Describe how to use the shift and logical operations to extract the day field from the packed date record in question 28. That is, wind up with a 16-bit integer value in the range 0..31.

30) Suppose you have a value in the range 0..9. Explain how you could convert it to an ASCII character using the basic logical operations.
31) The following C++ function locates the first set bit in the BitMap parameter starting at bit position start and working up to the H.O. bit. If no such bit exists, it returns -1. Explain, in detail, how this function works.

```cpp
int FindFirstSet(unsigned BitMap, unsigned start)
{
    unsigned Mask = (1 << start);
    while (Mask)
    {
        if (BitMap & Mask) return start;
        ++start;
        Mask <<= 1;
    }
    return -1;
}
```

32) The C++ programming language does not specify how many bits there are in an unsigned integer. Explain why the code above will work regardless of the number of bits in an unsigned integer.

33) The following C++ function is the complement to the function in the questions above. It locates the first zero bit in the BitMap parameter. Explain, in detail, how it accomplishes this.

```cpp
int FindFirstClr(unsigned BitMap, unsigned start)
{
    return FindFirstSet(~BitMap, start);
}
```

34) The following two functions set or clear (respectively) a particular bit and return the new result. Explain, in detail, how these functions operate.

```cpp
unsigned SetBit(unsigned BitMap, unsigned position)
{
    return BitMap | (1 << position);
}
```

```cpp
unsigned ClrBit(unsigned BitMap, unsigned position)
{
    return BitMap & ~(1 << position);
}
```

35) In code appearing in the questions above, explain what happens if the start and position parameters contain a value greater than or equal to the number of bits in an unsigned integer.
1.15 Programming Projects

The following programming projects assume you are using C, C++, Turbo Pascal, Borland Pascal, Delphi, or some other programming language that supports bitwise logical operations. Note that C and C++ use the “&”, “|”, and “^” operators for logical AND, OR, and XOR, respectively. The Borland Pascal products let you use the “and”, “or”, and “xor” operators on integers to perform bitwise logical operations. The following projects all expect you to use these logical operators. There are other solutions to these problems that do not involve the use of logical operations, do not employ such a solution. The purpose of these exercises is to introduce you to the logical operations available in high level languages. Be sure to check with your instructor to determine which language you are to use.

The following descriptions typically describe functions you are to write. However, you will need to write a main program to call and test each of the functions you write as part of the assignment.

1) Write to functions, `toupper` and `tolower`, that take a single character as their parameter and convert this character to upper case (if it was lowercase) or to lowercase (if it was uppercase) respectively. Use the logical operations to do the conversion. Pascal users may need to use the `chr()` and `ord()` functions to successfully complete this assignment.

2) Write a function “CharToInt” that you pass a string of characters and it returns the corresponding integer value. Do not use a built-in library routine like `atoi` (C) or `strtoInt` (Pascal) to do this conversion. You are to process each character passed in the input string, convert it from a character to an integer using the logical operations, and accumulate the result until you reach the end of the string. An easy algorithm for this task is to multiply the accumulated result by 10 and then add in the next digit. Repeat this until you reach the end of the string. Pascal users will probably need to use the `ord()` function in this assignment.

3) Write a `ToDate` function that accepts three parameters, a month, day, and year value. This function should return the 16-bit packed date value using the format given in this chapter (see “Bit Fields and Packed Data” on page 28). Write three corresponding functions `ExtractMonth`, `ExtractDay`, and `ExtractYear` that expect a 16-bit date value and return the corresponding month, day, or year value. The `ToDate` function should automatically convert dates in the range 1900-1999 to the range 0..99.

4) Write a “CntBits” function that counts the number of one bits in a 16-bit integer value. Do not use any built-in functions in your language’s library to count these bits for you.

5) Write a “TestBit” function. This function requires two 16-bit integer parameters. The first parameter is a 16-bit value to test; the second parameter is a value in the range 0..15 describing which bit to test. The function should return true if the corresponding bit contains a one, the function should return false if that bit position contains a zero. The function should always return false if the second parameter holds a value outside the range 0..15.

6) Pascal and C/C++ provide shift left and shift right operators (SHL/SHR in Pascal, “<<” and “>>” in C/C++). However, they do not provide rotate right and rotate left operators. Write a pair of functions, ROL and ROR, that perform the rotate tasks. Hint: use the function from exercise five to test the H.O. bit. Then use the corresponding shift operation and the logical OR operation to perform the rotate.